
Formalizing the Theory Concept in Nuprl

Jason Hickey
March 15, 1994

Outline

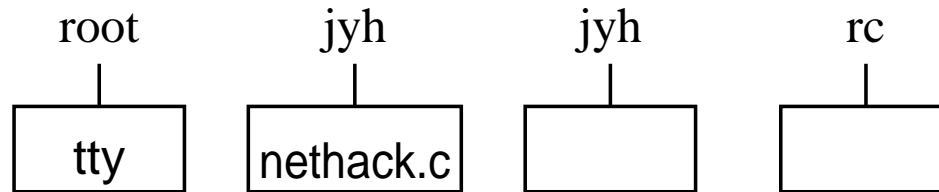
Goals of formalization

Goal generfization

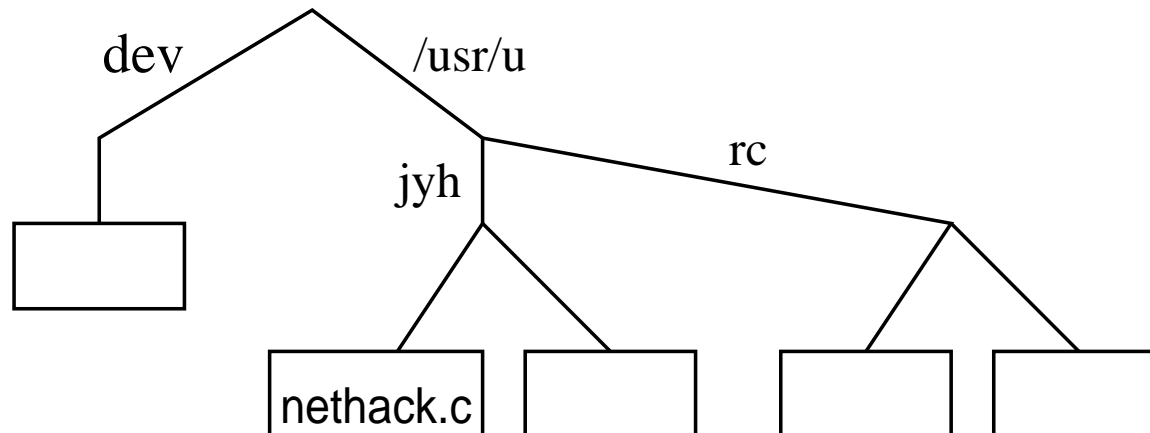
- Operate within a context of Theorems, Axioms, Definitions
- Primary goal is to extend the context
- Problems
 - flat name space
 - unrestricted access
 - generfl form of inheritance

Flat name space

Flat namespace



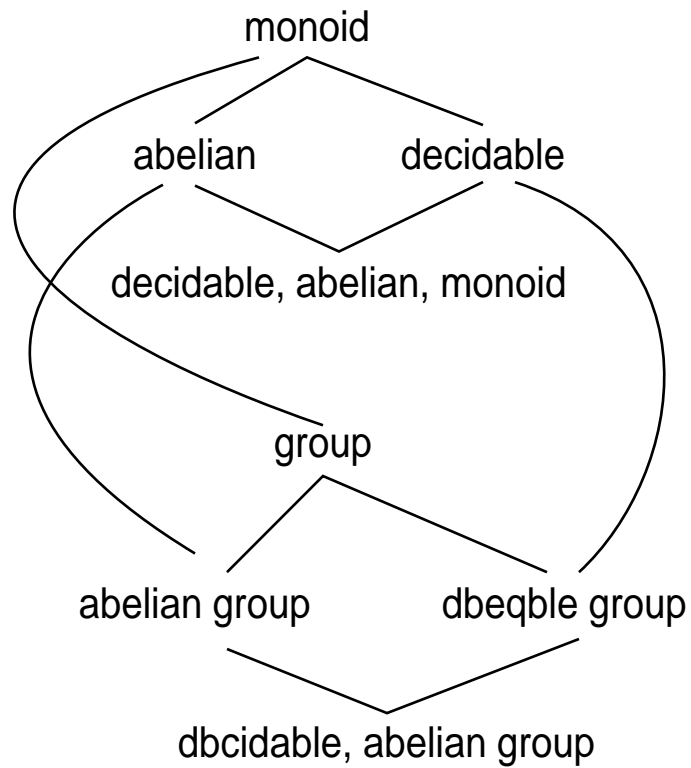
Hierarchy



Unrestricted Access

Inheritance

Running example



Reflection

- Want to reason about
 - proofs
 - theories
 - tactics

Object Orientation

- Abstract Data Typing
- Inheritance

* ABSmonoid

Monoid $fi\ g ==$ car: $Ufi\ g$

uni t: car

\pounds op: (car $!$ car $!$ car)

\pounds eq: (car $!$ car $!$ $Pfi\ g$)

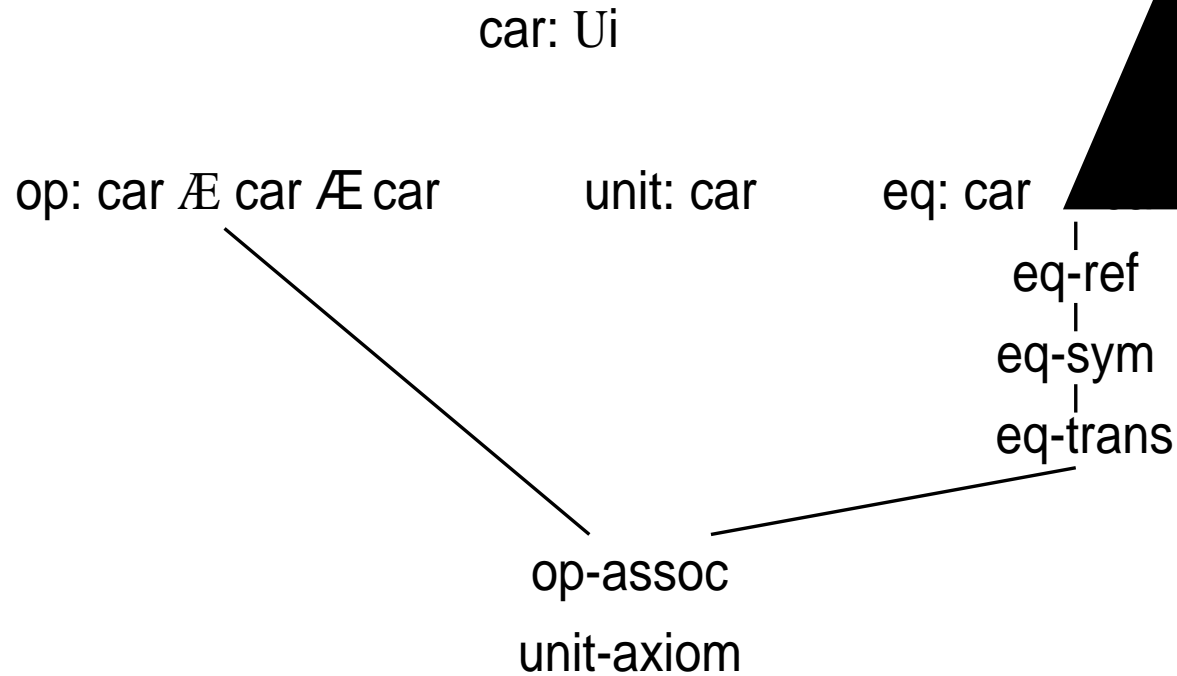
\pounds eq-ref: (\pounds a: car. a eq a)

\pounds eq-sym: (

Generalize the dependent product

- Number of elements is arbitrary

Least constraining ordering



TheoryItem

- Reference this axiom by name
- Use inhabitants of previous axioms by looking through ancestors
- What is an axiom?

* ABS theory

```
Theory fi g == rec(Theory.name: IDList
  £ label: Atom
  £ flags: List of Z
  £ lib: List of TheoryLibType
  £ preds: List of Theory
  £ Axiom?)
```

Axiom type

- Define lookup function
 - $\text{Ld0up}(\text{Type})$ name in preds

Axiom: $(\text{name: ID} \rightarrow$

\rightarrow

Ui

L

Example

- MonoidCar: <"car", [], lparents, lookup. Ui>

MonoidUnit:

<"unit",

[MonoidCar],

λ parents, lookup.

Lookup "car" in parents using MonoidCar>q7

Real Example

* ABS monoid car

MonoidCar $\models gfcar(g)$

* ABS monoid car thyMonoidCarThyfi g== Axiom:

Name = car/ ,

Flags = CNil ,

Lib = CNil ,

Preds = CNil ,

Axiom(par, pre) = MonoidCar

MonoidUnit

* ABS monoid unit preds

MonoidUnitPreds *fi g* == MonoidCarThy *fi g* : CNiI

* ABS monoid unit

MonoidUnit(parents, preds) == TheoryEnv[parents, preds]
car = car
in *fi g*
car

* ABS monoid unit thy

MonoidUnitThy *fi g* == Axiom:

Name = unit / ,
Flags = CNiI ,

MonoidOpAssoc

* ABS monoid unit preds

ds $fi\ g$ $fi\ g : CNil$

* ABS monoid unit

parents, preds) == TheoryEnv[parents, preds]

car = car

in $fi\ g$

car

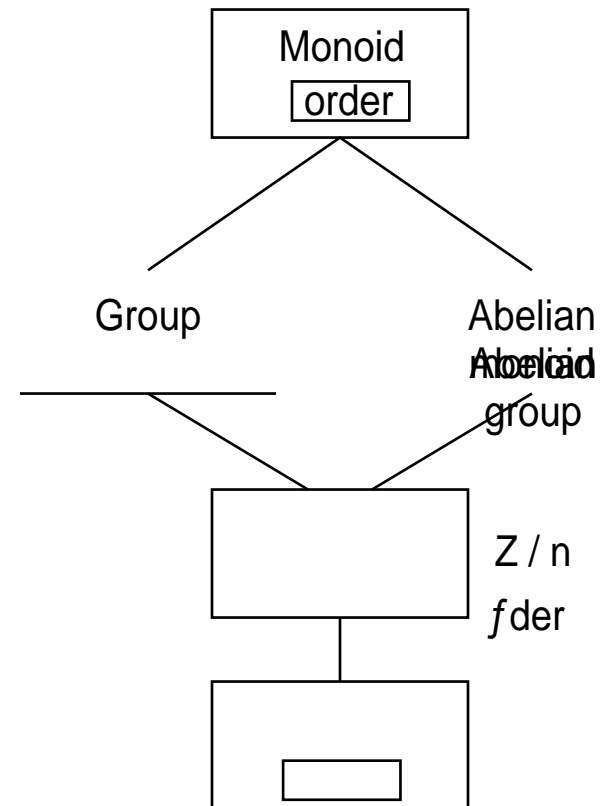
* ABS monoid unit thy

M1idUnitThy $fi\ g$ == Axiom:

Name = unit/

Theorems

- Axiom: "every monoid has an order"
- Theorem extract: function that computes the order
- May want several proofs of the theorem depending on the particular monoid



Why do we want Theorems?

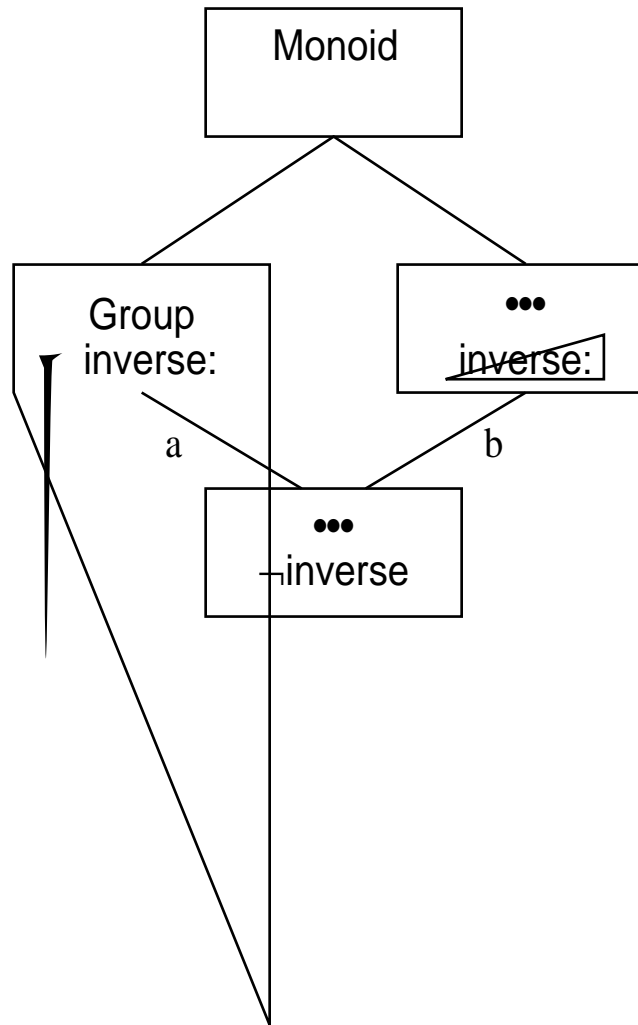
- To show something is true
 - "Every monoid of prime order is cyclic"
- To compute
 - "Every monoid has an order"
- Don't handle extracts!

A Theorem is just a link from an axiom to an object in the

Naming considerations

Can pull in any theory

Must inherit all axioms as well



Problems

- `rec(Theory. ●●●)` does not work!
 - Typing rule is too simple.
- Build type on top of Y-combinator

Y-combinator

- Can hide all instances of bounds by using derived rules
- All proofs are by (normal) induction
- No fixed-point semantics

Data structure

- In a functional language where equality is undecidable, how can we represent a DAG? (Would like an equality over strings)
- My solution
 - need an arbitrary total order.
 - the DAG are integer indices into the total order
 - Not elegant!

Results

- Can read about theories

A theory is a normal object, reading is typically by induction

- Examples

- Knowledge monotonically increases
- Equivalence of theories
- "Lifting" of theorems: any theorem can be lifted to be beneath its immediate predecessors

Results II

- Naming is better

Problems

Future work

- Package theories in a hierarchy
 - Conceptual blocks
Assist in naming

TheoryType

```

* ML theory type switch ml
% Given a 'TheoryItem' type fig(n)', construct a product of all the axioms
  in the theory. This effectively all the junk out of a theory
  tdV4(provide)-501(a)-507(type)-516(that)-516(can)-537(be)-510(quantified)-510(over.)]TJ-1.032 -1.2 TD
  'case
  of inl ( ) =>
    parents: TheoryTypech(inr      , item.preds)
     $\mathcal{E}$  Switch x = item.term of
      Axiom ! x parents (, name.Lookup(preds, f name in parents using
      Theorem ! Unit ! Unit | inr ( ) =>
        Case item of
        CNil => Unit
        hd::tl => TheoryTypeSwitch(      , hd.theory)  $\mathcal{E}$  TheoryTypeSwitch(

```

TheoryLookup

```
* ML theory lookup switch ml
% Given a 'TheoryItemType fi g(n)', and a term of 'TheoryType(item)',
  and a name of type 'ID', return the term inhabiting the axiom by that name.
%
add rec def 'Lookup[switch] name in term using item'
  'case switch
    of inl ( ) =>
      if name 2 item.name
      then inr term.2
      else wi Lookup[inr] name in term.1 using item.preds
    fi
  inr ( ) =>
    Case item of
    CNil => inl
    hd::tl => case wi Lookup[inl] name in term.1 using hd.theory
      of inl ( ) =>
        Lookup[inr i] name in term. 2 using tl
      inr ( x) =>
        inr x;
```

```

* ML theory lookup type switch ml
% Give a 'TheoryItem type fig(n)', and a term of 'TheoryType(item)',
  and a name 'ID', return the axiom of that name.
%
add rec def 'Lookup(type)[switch] name in term using item'

of case (w) then
  if name = item.name
  then inl (term)
  else if ! (Lookup(type)[inr] name in term.1 using item.preds)
  then ! (Lookup(type)[inr] name in term.1 using item.preds)
  fi
  | inr() =>
    Case item of
    CNil => inl Unit
    hd::tl => case Lookup(type)[inl] name in term.1 using hd.theory
              of inl() =>
                  Lookup(type)[inr] name in term.2 using tl

```