
Abstract Programming in Nuprl

“The Majority Vote Algorithm”

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Outline

- **Abstraction**
- **Majority vote algorithm**
- **Abstract implementation**
- **Solution to majority vote**
- **Concrete implementation**
- **Compare with other languages**

Themes of high-level programming

- **Abstraction**
 - Hides underlying implementation
 - Packages structure into a modules for reuse

Transformation-real abstraction

- We should be able to specify an abstraction without an implementation
- Virtual program is transformed to an executable program by replacing the virtual instruction with a concrete implementation.

Remove all duplic

- **Given a bag X of size n and the majority k**

- If b has at least k distinct values, delete them
- Perform at most $n \div k$ times

Every value occurring more than $n/2 \div k$ in X is in X

Delete k distinct values from X

Multiset

- **Multiset is the major data type–target for abstraction**
- **Follow presentation given in Manna and Waldinger**
- **Basic elements and operations**
 - empty set
 -
 -

Multiset Signature

car: Ufi g

- £ empty: car
- £ gen: (car ! Atom ! car)
- £ member: (car ! Atom ! Pfi g)
- £ equal: (car ! car ! Pfi g)
- £ syneq: (car ! car ! Pfi g)
- £ atomeq: (Atom ! Atom ! Pfi g)
- £ ind: MSetIndTypefgAtom, car, empty, gen, syneq)
- £ mem p: MSetMemType(Atom, car, empty, member, gen, equal, atomeq)
- £ unique p: MSetUniqType(Atom, car, empty, gen, equal, syneq, atomeq)
- £ equality p: MSetEqType(Atom, car, gen, equal)
- £ atomeq p: MSetAEqType(Atom, atomeq)
- £ syneq p: MSetSeqType(car, syneq)
- £ B

Multiset Axioms

- **induction principle**
- **uniqueness (free generation)**
- **equality axioms**
- **membership axioms**

Multiset Theorems

* THM mset ind thm

\mathcal{A} Atom: Ufi g

\mathcal{M} : MSetSi gfi g (Atom)

\mathcal{P} : m. car ! Pfi g

(\mathcal{A} . m.empty \mathcal{M}) P[b]

) (\mathcal{A} b: m. car. \mathcal{A}) P[b [m x]]

) f \mathcal{A} b: m. car. P[b] g

* THM mset mem empty

\mathcal{A} Atom: Ufi g . \mathcal{M} : MSetSi gfi g (Atom). \mathcal{A} x: Atom. : (x \mathcal{M} m.empty)

* THM mset mem thm

\mathcal{A} Atom: Ufi g

\mathcal{M} : MSetSi g



Other functions and Theorems

- **Decomposition**
- **Equality Theorems**
- **MSet union**
- **MSet difference**
- **etc.**

Return to Majority Vote

- **Define a predicate that is true when an element has the majority vote:**
 - $\text{MajVoteProp}[k](\text{Atom}, m, b, x)$
 - x has a majority in b
- **Define a predicate that is true when a multiset is a set of k distinct elements**
 - $\text{MajDistinctProp}[k](\text{Atom}, m, b)$
 - b is a set, and it has k distinct elements
- **Define a k -step reduction**
 - $\text{ReducedBag}[k](A, m, b \rightarrow c)$
 - There is a set d with k distinct elements and $c \gg d = b$

Majority Vote II

- **Define a predicate fdPr a reduced bag**
 - $\text{MajRedBag}[k](\text{Atom}, m, b \rightarrow c)$
 - c has fewer than k distinct elements
 - all elements x having a majority vote in b have a majority vote in c

The dms

- The elements with a majority vote are in every reduced bag
- Normalizat
 - if $b - c = d$ and $d \rightarrow e$ and c is E st of k distinct elements, then $b \rightarrow e$
- Main Result:
 - Fd• every bag, there is a reduced bag
 - Proof is by induction on the siz of the bag
 - If Ethe bag has fewer than k distinct elements, done
 - Otherwise, •move k distinct elements. The Eresult Eholds by induction.

Comparison of the two developments

- **Gries logic separates computation from verification**
 - Logic is classical
 - Every proof is for a particular program
- - Better--why do more than necessary?
Worse--computation is hidden

ThorVgorithm is accompanied by a proof

Transformation

- It would be nice to try out this algorithm
- **Brief presentation of an implementation based on lists.**
 - empty: []
 - add: cons
 - member: member

Transformation II

- All the properties defined in the previous slide
- Other options:
 - function that returns the multiplicity of an element in a bag
- **Computation of majority vector**
 - $\text{ReducedBag}([1;1;2]) = [1]$

What I learned in this process

- **Time to learn the system--its huge!**
 - **Existing documentation is good, but its incomplete**
 - **Power of abstraction**
 - **Power of transformation**
 - **Delayed/lazy development**
- All these things are poss**

Proofs

- **These proofs are tedious!**
Many of the proofs are small, and conceptually

- " b: Bag " x: Atom b » x p 0
- Three of four steps to prove this

Many such trivial proofs

- **Many well-formulated proofs**
Many well-formulated proofs set of rules

Incremental Development

- **Small modifications wipe out an entire theorem or library**
 - Save the proof tree?
- **Hypothesis numbering**
- **Libraries as first class objects**
 - Axioms as assumptions
 - Naming

Advantages

- **Higher level of abstraction than any other language I know**
- **Transformation is natural**
- **The problem with tedious proofs can be solved with appropriate ATP's**

Other languages

- **Most languages do not have specifications**

Example in C++

```
class MSet f
private:
    LinkedList car;
public:
    Bool aeq(int; int);
    Bool eq(MSet x);
    Bool member(Atom x);
    void add(Atom x);
    Atom head
    Atom tail
    MSet extract(Atom xg;
```


Features of Object Orientation

Inheritance II

Modify features

Conclusion

- **Powerful abstraction and transformation mechanism**
- **Majority vote algorithm has a simple solution**
- **Proofs sn be tedious, but san be fixed with tastis help**
- **We can provide some features of an object oriented language, do we want them all?**
- **What is the sorrect model for a theory?**